#### **Schema and Data Translation**

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Based on

Paper in EDBT 1996 (with R. Torlone) Paper in EDBT 2006 (with P. Cappellari and P. Bernstein) Demo in Sigmod 2007 (with P. Cappellari and G. Gianforme)

InfInt -- Bertinoro, October 4, 2007

## **Terminology: a warning**

Model Mgmt people	Traditional DB people
Meta-metamodel	Metamodel
Metamodel	Model
Model	Schema

#### Schema and data translation

- Schema translation:
  - given schema S1 in model M1 and model M2
  - find a schema S2 in M2 that "corresponds" to S1
- Schema and data translation:
  - given also a database D1 for S1
  - find also a database D2 for S2 that "contains the same data" to D1

## A long standing issue

- Translations from a model to another have been studied since the 1970's
- Whenever a new model is defined, techniques and tools to generate translations are studied
- However, proposals and solutions are usually model specific:
  - Given an ER schema, find the suitable relational schema that "implements" it
    - the original paper (Chen 1976) contains the basics
    - further discussions by many (e.g. Markowitz and Shoshani 1989)
    - illustrated in every textbook
  - Similarly with
    - any other conceptual model and any other logical one
    - XML and relational (or object)

#### We have been doing this for a while

- Initial work more than ten years ago (Atzeni & Torlone, 1996)
- Major novelty recently (Atzeni, Cappellari & Bernstein, 2006; Atzeni, Cappellari & Gianforme, 2007)
  - translation of both schemas and data
  - data-level translations generated automatically, from schema-level ones

## A simple example

- An object relational database, to be translated in a relational one
- Source: the OR-model
- Target: the relational model



## Example, 2





- Does the OR model allow for keys?
- Assume EmpNo and Name are keys

Er	nployees		]			
<u>EmpNo</u>	Name	Dept		[	De	epartments
134	Smith	A		Name	e	Address
201	Jones	В		A	Γ	5, Pine St
255	Black	A		В	/	10, Walnut St
302	Brown	null		$\bigcirc$		

## Example, 3

	Employees					
	EmpNo	Name	Dept			
E#1	134	Smith	D#1			
E#2	201	Jones	D#2			
E#3	255	Black	D#1			
E#4	302	Brown	null			



- Does the OR model allow for keys?
- Assume no keys are specified

$\square$	Employ	yees				
EmpID	EmpNo	Name	Dept	$\frown$	Depart	ments
1	134	Smith	1	DeptID	Name	Address
2	201	Jones	2	1	А	5, Pine St
3	255	Black		2	В	10, Walnut St
4	302	Brown	null			
$\overline{}$						

## Many different models (and variations ...)



## **Heterogeneity**

- We need to handle artifacts and data in various models
  - Data are defined wrt to schemas
  - Schemas are defined wrt to models
  - How models can be defined?



## A metamodel approach

- The constructs in the various models are rather similar:
  - can be classified into a few categories (Hull & King 1986):
    - Lexical: set of printable values (domain)
    - Abstract (entity, class, ...)
    - Aggregation: a construction based on (subsets of) cartesian products (relationship, table)
    - Function (attribute, property)
    - Hierarchies
    - ...
- We can fix a set of metaconstructs (each with variants):
  - lexical, abstract, aggregation, function, ...
  - the set can be extended if needed, but this will not be frequent
- A model is defined in terms of the metaconstructs it uses

#### The metamodel approach, example

- The ER model:
  - Abstract (called Entity)
  - Function from Abstract to Lexical (Attribute)
  - Aggregation of abstracts (Relationship)
  - ...

. . .

- The OR model:
  - Abstract (Table with ID)
  - Function from Abstract to Lexical (value-based Attribute)
  - Function from Abstract to Abstract (reference Attribute)
  - Aggregation of lexicals (value-based Table)
  - Component of Aggregation of Lexicals (Column)

## The supermodel

- A model that includes all the meta-constructs (in their most general forms)
  - Each model is subsumed by the supermodel (modulo construct renaming)
  - Each schema for any model is also a schema for the supermodel (modulo construct renaming)

#### The metamodel approach, translations

- The constructs in the various models are rather similar:
  - can be classified into a few categories ("metaconstructs")
  - translations can be defined on metaconstructs,
    - and there are "standard", accepted ways to deal with translations of metaconstructs
    - they can be performed within the supermodel
  - each translation from the supermodel SM to a target model
     M is also a translation from any other model to M:
    - given n models, we need n translations, not n<sup>2</sup>

#### **Generic translation environment**



#### **Translations within the supermodel**

- We still have too many models:
  - Just within simple ER model versions, we have 4 or 5 constructs, and each has several independent features which give rise to variants
    - for example, relationships can be
      - binary or N-ary
      - with all possible cardinalities or without many-tomany
      - with or without the possibility of specifying optionality
      - with or without attributes

— ...

- Combining all these, we get hundreds of models!
- The management of a specific translation for each model would be hopeless

## **Translations, the approach**

- Elementary translation steps to be combined
- Each translation step handles a supermodel construct (or a feature thereof) "to be eliminated" or "transformed"
- A translation is the concatenation of elementary translation steps



#### **Complex translations**



## **Translations**

- Basic translations are written in a variant of Datalog, with OID invention
  - We specify them at the schema level
  - The tool "translates them down" to the data level
  - Some completion or tuning may be needed

## **A Multi-Level Dictionary**

- Handles models, schemas and data
- Has both a model specific and a model independent component
- Relational implementation, so Datalog rules can be easily specified



mM 🔇	mSM
М	SM
i-M	i-SM

## The supermodel description

MSM-Construct					
OID	Name	IsLex			
1	AggregationOfLexicals	F			
2	ComponentOfAggrOfLex	Т			
3	Abstract	F			
4	AttributeOfAbstract	Т			
5	F				

MSM-Reference					
<u>OID</u>	Name	Со	nstru	uct	Target
30	Aggregation		2		1
31	Abstract		4		3
32	Abstract1		5		3
33	Abstract2		5		3
			$\bigcirc$		

MSM-Property					
<u>OID</u>	Name	С	Constr.		Туре
11	Name		1		String
12	Name		2		String
13	lsKey		2		Boolean
14	IsNullable		2		Boolean
15	Туре		2		String
16	Name		3		String
17	Name		4		String
18	Isldentifier		4		Boolean
19	IsNullable		4		Boolean
20	Туре		4		String
21	IsFunct1		5		Boolean
22	IsOptional1		5		Boolean
23	Role1		5		String
24	IsFunct2		5		Boolean
25	IsOptional2		5		Boolean
26	Role2		5		String

$\left( \right)$	mМ	mSM
	М	SM
	i-M	i-SM

#### **Model descriptions**

![](_page_23_Figure_2.jpeg)

. . .

	MSM-Construct						
	OID	IsLex					
	1	AggregationOfLexicals	F				
1	2	ComponentOfAggrOfLex	Т				
	3	Abstract	F				
	4	AttributeOfAbstract	Т				
	5	BinaryAggregationOfAbstracts	F				

	MSM-Property					
<u>OID</u>	OID Name Construct Type					

MSM-Reference					
OID Name Construct					

P. Atzeni

...

. . .

. . .

	mМ	mSM
	М	SM
	i-M	i-SM

## Schemas in a model

![](_page_24_Figure_2.jpeg)

![](_page_25_Picture_0.jpeg)

![](_page_25_Figure_1.jpeg)

mSM

SM

i-SM

mΜ

Μ

i-M

M SM i-M i-SM Instances in the supermodel									
SM-AttributeOfAbstract									
	ØD	Setre fab	s <b>OND</b> ame In	st <b>QlD</b> en	t i <b>¥alut</b> able	TypieAb <b>≴03b</b> ΩrOID			
	4000	1 401	EmpNo	1 T	754372	Int 3010 301			
3817 301 Departments	4000	402	Name	1 F	John Boe	Text 3010 301			
201 3 Clerks	404		Name	Т	F	Char 302			
3010 301 Offices	4005	1 402	Address	1 F	Bob Whaite	Text 3011 302			
	501	3	Code	T	F	Int 201			
Supermodel instances									

![](_page_27_Figure_0.jpeg)

## **Translations**

- Basic translations are written in a variant of Datalog, with OID invention
  - We specify them at the schema level
  - The tool "translates them down" to the data level
  - Some completion or tuning may be needed

## A basic translation

- From (a simple) binary ER model to the relational model
  - a table for each entity
  - a column (in the table for E) for each attribute of an entity E
  - for each M:N relationship
    - a table for the relationship
    - columns ...
  - for each 1:N and 1:1 relationship:
    - a column for each attribute of the identifier ...

#### A basic translation application

![](_page_30_Figure_1.jpeg)

	Employees							
Em	<u>pNo</u>	Name		Affiliation				
	Departments							
	Na	me Address		lress				

## A basic translation (in supermodel terms)

- From (a simple) binary ER model to the relational model
  - artable for each abstract
  - a cohuprom (in the trade leggore E) from efacte attribute boft a rofe altistr Ect
  - for each M:N æggtiegnsthipp of abstracts ...
    - a table for the relationship
    - columns ...
  - for each 1:N and 1:1 relationship:
    - a column for each attribute of the identifier ...

## "An aggregation of lexicals for each abstract"

```
SM_AggregationOfLexicals(
   OID: #aggregationOID_1(OID),
   Name: name)
\leftarrow
SM_Abstract (
   OID: OID,
   Name: name );
```

## **Datalog with OID invention**

- Datalog (informally):
  - a logic programming language with no function symbols and predicates that correspond to relations in a database
  - we use a non-positional notation
- Datalog with OID invention:
  - an extension of Datalog that uses Skolem functions to generate new identifiers when needed
- Skolem functions:
  - injective functions that generate "new" values (value that do not appear anywhere else); so different Skolem functions have disjoint ranges

#### "An aggregation of lexicals for each abstract"

![](_page_34_Figure_1.jpeg)

- the value for the attribute Name is copied (by using variable n)
- the value for OID is "invented": a new value for the function #aggregationOID\_1(OID) for each different value of OID, so a different value for each value of SM\_Abstract.OID
- Skolem functions are materialized in the dictionary:
  - Represent the mapping

## "An aggregation of lexicals for each abstract"

![](_page_35_Figure_1.jpeg)

# "A component of the aggregation for each attribute of abstract"

![](_page_36_Figure_1.jpeg)

- Skolem functions
  - are functions
  - are injective
  - have disjoint ranges
- the first function "generates" a new value
- the second "reuses" the value generated by the first rule

![](_page_37_Figure_0.jpeg)

#### **Generating data-level translations**

- Same environment
- Same language
- High level translation specification

Supermodel description (mSM)

![](_page_38_Figure_5.jpeg)

## **Translation rules, data level**

![](_page_39_Figure_1.jpeg)

## Correctness

- Usually modelled in terms of information capacity equivalence/dominance (Hull 1986, Miller 1993, 1994)
- Mainly negative results in practical settings that are non-trivial
- Probably hopeless to have correctness in general
- We follow an "axiomatic" approach:
  - We have to verify the correctness of the basic translations, and then infer that of complex ones

## **Experiments**

- A significant set of models
  - ER (in many variants and extensions)
  - Relational
  - OR
  - XSD
  - UML
- Demo

![](_page_42_Figure_0.jpeg)

![](_page_43_Figure_0.jpeg)

## Summary

- ModelGen was studied a few years ago
- New interest on it within the "Model management" framework
- New approach
  - Translation of schema and data
  - Visible (and in part self generated) dictionary
  - Visible and modifiable rules
  - Skolem functions describe mappings

#### Issues

- ModelGen in the model management scenario:
  - Round-trip engineering (and more)
- Customization of translations
- Off-line translation vs run-time
- "Correctness" of data translation wrt schema translation:
  - compare with data exchange
- Schematic heterogeneity and semantic Web framework:
  - What if the distinction between schemas and instances is blurred?
- Materialized Skolem function & provenance

#### Modelgen in model management

- Round trip engineering (Bernstein, CIDR 2003)
  - A specification (for example ER or UML) and an implementation (for example, relational)
  - then a change to the implementation: want to revise the specification
- We need a translation from the implementation model to the specification one

![](_page_46_Figure_5.jpeg)

#### **Round trip engineering**

![](_page_47_Figure_1.jpeg)

m2 = Match (I1,I2) m3 = Compose (m1,m2) I2'= Diff(I2,m3) <S2',m4 > = Modelgen(I2') ... Match, Merge

#### Another problem in the picture: data exchange

- Given a source S1 and a target schema S2 (in different models or even in the same one), find a translation, that is, a function that given a database D1 for S1 produces a database D2 for S2 that "correspond" to D1
- Often emphasized with reference to materialized solutions

## Integration

Given two or more sources, build an integrated schema or database

![](_page_49_Figure_2.jpeg)

#### **Schema translation**

• Given a schema find another one with respect to some specific goal (another model, better quality, ...)

![](_page_50_Figure_2.jpeg)

#### Data exchange

• Given a source and a target schema, find a transformation from the former to the latter

![](_page_51_Figure_2.jpeg)

#### Schema translation and data exchange

- Can be seen as complementary:
  - Data translation = schema translation + data exchange
    - Given a source schema and database
    - Schema translation produces the target schema
    - Data exchange generates the target database
- In model management terms we could write
  - Schema translation:
    - <S2, map12> = ModelGen (S1,mod2)
  - Data exchange:
    - i2 = DataGen (S1,i1,S2,map12)